

# Mathematics in Image Processing

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- We shall discuss this only in the discrete setting

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**what are the best questions to ask??**
- Two issues: (i) Enough information in  $y$  to determine  $x$ ;  
(ii) How to extract the information  $y$  holds about  $x$ :  
**Decoder**

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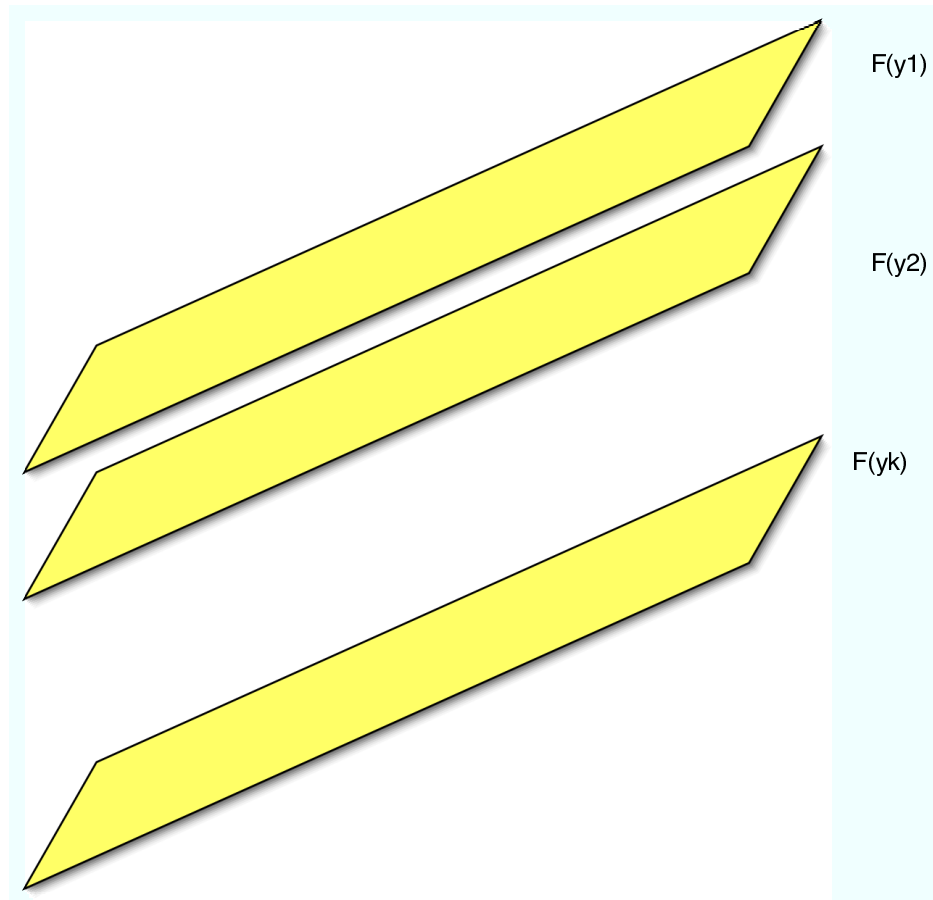
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# The sets $\mathcal{F}(y)$



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- $\bar{x} := \Delta(\Phi(x))$  is our approximation to  $x$  from the information extracted
- Note that all  $x \in \mathcal{F}(y)$  are approximated by the same  $\bar{x}$

# Measuring Sparsity

- Compressed Sensing models signals as sparse in some basis
- By linear transformation, we can assume **WOLOG** that  $x$  is sparse with respect to the canonical basis on  $\mathbb{R}^N$
- The support of  $x$  is  $\text{supp}(x) := \{i : x_i \neq 0\}$
- $\Sigma_k := \{x : \#\text{supp}(x) \leq k\}$
- Note that  $\Sigma_k$  is a union of  $k$  dimensional subspaces:  
 $\Sigma_k = \cup_{\#(T)=k} X_T$  where  $X_T = \{x : \text{supp}(x) \subset T\}$
- **First Question:** Given  $k, N$  what is the smallest  $n$  for which there is  $(\Phi, \Delta)$  such each vector in  $\Sigma_k$  is captured exactly  $\Delta(\Phi(x)) = x, \quad x \in \Sigma_k$

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- Answer  $n = 2k$

# What matrices do the job?

- $\Phi = [v_1, \dots, v_N]$ ,  $v_1, \dots, v_N$  columns of  $\Phi$
- We say  $\Phi$  has the independence property (IP) of order  $k$  if all choices of  $k$  column vectors are independent
- If  $T = \{i_1, \dots, i_m\}$  is a set of column indices
- $\Phi_T = [v_{i_1}, \dots, v_{i_m}]$  is the  $n \times \#(T)$  submatrix of  $\Phi$  formed from these columns
- IP means  $\Phi_T^* \Phi_T := (\langle v_i, v_j \rangle)_{i,j \in T}$  is invertible (positive eigenvalues) whenever  $\#(T) = k$

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**Theorem:** If  $\Phi$  is any  $n \times N$  matrix and  $2k \leq n$ , then the following are equivalent:

- There is a  $\Delta$  such that  $\Delta(\Phi(x)) = x$ , for all  $x \in \Sigma_k$ ,
- $\Sigma_{2k} \cap \mathcal{N}(\Phi) = \{0\}$ ,
- the matrix  $\Phi_T$  has the independence property of order  $2k$ .

# Optimal Matrices

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- Vandermonde matrix. Choose  $x_1 < x_2 < \dots < x_N$

$$\begin{pmatrix} 1 & 1 & \dots & 1 \\ x_1 & x_2 & \dots & x_N \\ \vdots & \vdots & \dots & \vdots \\ x_1^{2k-1} & x_2^{2k-1} & \dots & x_N^{2k-1} \end{pmatrix} \quad (3)$$

# Naive Decoding

$$\Delta(y) := \underset{z \in \Sigma_k}{\text{Argmin}} \|y - \Phi(z)\|_{\ell_2^n}$$

- $X_T := \{z : \text{supp}(z) \subset T\}$
- $x_T := \underset{z \in X_T}{\text{Argmin}} \|y - \Phi z\|_{\ell_2^n} \rightarrow x_T = [\Phi_T^* \Phi_T]^{-1} \Phi_T y$
- $T^* := \underset{\#(T)=k}{\text{Argmin}} \|y - \Phi(x_T)\|_{\ell_2^n}$
- $\Delta(y) := x_{T^*}$

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- We would need this norm controlled for any  $T$  of size  $k$

# Optimal Stable Systems

- Candes-Romberg-Tao; Donoho: Compressed Sensing
- Two important discoveries: good matrices, decoding
- Restricted Isometry Property (RIP) of order  $k$ : There exists  $0 < \delta = \delta_k < 1$  such that

$$(1 - \delta)\|x\|_{\ell_2^N}^2 \leq \|\Phi(x)\|_{\ell_2^n}^2 \leq (1 + \delta)\|x\|_{\ell_2^N}^2, \quad x \in \Sigma_k$$

- Equivalently the eigenvalues of  $\Phi_T^* \Phi_T$  are in  $[1 - \delta, 1 + \delta]$
- Decode by  $\ell_1$  minimization

$$\Delta(y) := \inf_{x \in \mathcal{F}(y)} \|x\|_{\ell_1^N}$$

- Candes-Tao: If  $\Phi$  satisfies the RIP of order  $3k$  then given any  $x \in \Sigma_k$  we have  $\Delta(\Phi(x)) = x$  for the  $\ell_1$  minimization decoder. Moreover, the decoding is stable

# Building matrices

- How can we build matrices that satisfy RIP for the largest value of  $k$
- Given  $n, N$  we can construct such matrices for any  $k \leq c_0 n / \log(N/n)$
- The additional  $\log(N/n)$  is the price we pay for stability
- This is the largest possible range of  $k$
- How can we construct such  $\Phi$ ?
- We want to create a lot of vectors  $v_1, \dots, v_N$  in  $\mathbb{R}^n$  such that any choice of  $k$  of them are far from being linearly dependent

# Three constructions

- We choose at random  $N$  vectors from the unit sphere in  $\mathbb{R}^n$  and use these as the columns of  $\Phi$
- We choose each entry of  $\Phi$  independently from the Gaussian distribution with mean 0 and variance  $n^{-1}$
- We use Bernoulli process and create a matrix with entries  $\frac{\pm 1}{\sqrt{n}}$ , with equal probability
- With high probability each of these random constructions yields a matrix  $\Phi$  with RIP of order  $k$  for the (largest) range  $k \leq c_0 n / \log(N/n)$

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- Probability is only used to prove existence of  $\Phi$ . The sensing algorithm is constructive (not probabilistic) once we find a  $\Phi$ .

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- Park City, Utah: July 2009